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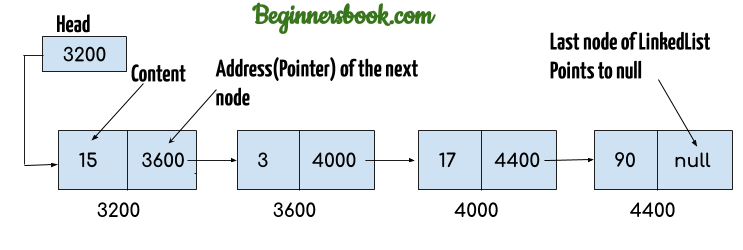
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# What Is Linked List ?

A Linked List can be defined as a collection of objects called nodes that are randomly stored in the memory.

* A node contains two fields :
* data stored at that particular address
* the pointer which contains the address of the next node in the memory.

****

# **Structure Of Linked List Vs Structure Of Array In The Memory**

Array Stored As One Block Or Contiguous Locations

| Address | Data |
| --- | --- |
|  |  |
|  |  |
| Beginning Of Array | |
| 0x01 | 23 |
| 0x02 | 5 |
| 0x03 | 47 |
| End Of Array | |
|  |  |
|  |  |

Linked List Stored in Different Or Random Locations And The Items Connected With Each Other By Next Pointers.

| Address | Data | Next |
| --- | --- | --- |
| 0x01 | 23 | 0x04 |
| 0x02 |  |  |
| 0x03 |  |  |
| 0x04 | 5 | 0x06 |
| 0x05 |  |  |
| 0x06 | 47 | null |
| 0x07 |  |  |

# Advantages Of Linked List :

## 1.Linked List Dynamic Structure Advantages :

**1.1 The Size Of Linked List is not Required When Linked List is Created At Compile Time.**

So Linked List Can Be Useful Over Array When The Size Of Elements Is Unknown And Depend On the inputs of user or data On the Run Time.

**1.2 The Size Of Linked List Is Dynamic At Run Time.**

The Size Of Array Is Fixed , So When An Array is Created , The Size Of Array Can Not Be Changed ( Extended Or Shrinked ).

Example : If You Declare Array With Size Of 200 Elements

- You Can not Change This Size At Run Time To 300 If You Need To Add New Elements.

- You Can not Change This Size At Run Time To 100 If You Need To Delete Unused Elements To Save Memory Space.

With Linked List in The Run Time :

* The Size Of Linked List Can Be Extended By Adding New Nodes Until We Reach The Max Allowed Size on The Memory.
* The Size Of Linked List Can Be Shrinked By Deleting Unwanted Nodes So we can save memory space.

## 

## 

## 2.Linked List Operations Advantages :

**insert or delete any element at any position is very fast.**

These two operations take a constant time (O) = 1.

And this is one of the biggest advantages of using the linked list. We will Explain All Operations In Details In The Next Sections.

## 3.Linked List Memory Advantages :

no unused elements can appear and waste memory space like arrays.

Example : If we need now to store data = { 1 , 3 , 5 } and maybe in the future we will add { 2 , 4 }.

**In array** if we declare the size of array = 5

| 1 | 3 | 5 | unused | unused |
| --- | --- | --- | --- | --- |

**linked list** Used Only Numbers Of Nodes As Required and When New Data Coming The Pointer Of Last Node Will Reference To It.

| Current address | data | Next node address |
| --- | --- | --- |
| 100 | 1 | 200 |
| The space between 100 and 200 used for other purposes during the execution of the program | | |
| 200 | 5 | 300 |
| The space between 200 and 300 used for other purposes during the execution of the program | | |
| 300 | 3 | null |

# Disadvantages of Linked List :

## 1.Linked List Operations Disadvantages :

1.1 Access Operations Are Very Slow [ (O) = N ] When Compared With Arrays Which Have [ (O) = 1 ].

linked list dose not have indexes Like Array so Random access is not allowed. We have to access elements sequentially starting from the first node Until We Reach The Required Node.Unlike The Array Which Can Reach To The Required Element And Access It In One Step Through The Equation Of The Base Address As We See in The Previous Chapter.

1.2 Search Operation is Slow [ (O) = N ] When Compared With Data Structure Like Binary Search Tree Which Has [ (O) = Log n ]

## 2.Linked List Memory Disadvantages

The Linked List Takes Extra Space Over The Array To Store The Pointers Of The Next Node Address.

Arrays store only the data of elements.

but linked list need to store the data and the address of next node.

# Types Of Linked List :

1- Singly Linked List

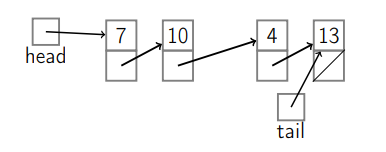
2- Doubly Linked List

3- Circular Linked List

# **Singly Linked List**

## What Is a Singly Linked List ?

Singly Linked List is a One way chain or can be traversed only in one direction. In other words, we can say that each node contains only the next pointer, therefore we can not traverse the list in the reverse direction.



## Singly Linked List Structure :

1- Header : Used To Point To The First Node In The Linked List

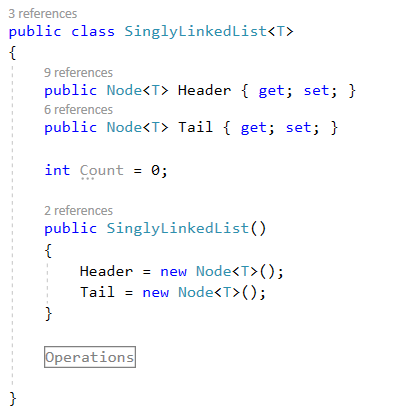
2- Tail : Used To Point To The Last Node In The Linked List

3- Data Nodes : The Other Nodes Between Header And Tail That used to Store Actual Data Of LinkedList and every node consist of two fields :

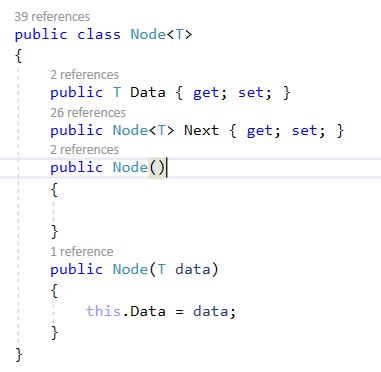
* Data field : To Store Data
* Next Field : To Store The Address Of The Next Node.

## Singly Linked List Structure Code :

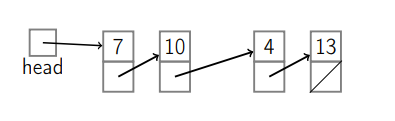
Singly Linked List Class :



Node Class :



## Operations In Singly Linked List



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## 1- Insert Operations

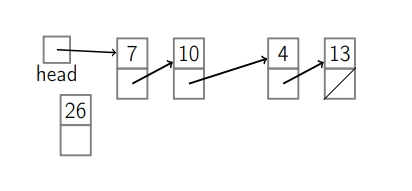
## 1.1 Insert At First

The new node will be the first node in the linked list.

Called Also in some reference : Insert After Header or Push Front

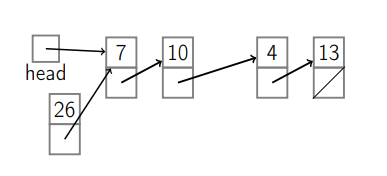
### - Steps :

1. Create New Node

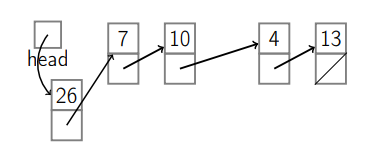


2. If LinkedList Is Not Empty

2.1 The Next Of The New Node Will Point To The Current First Node



2.2 The Next Of Header Will Point To The New Node

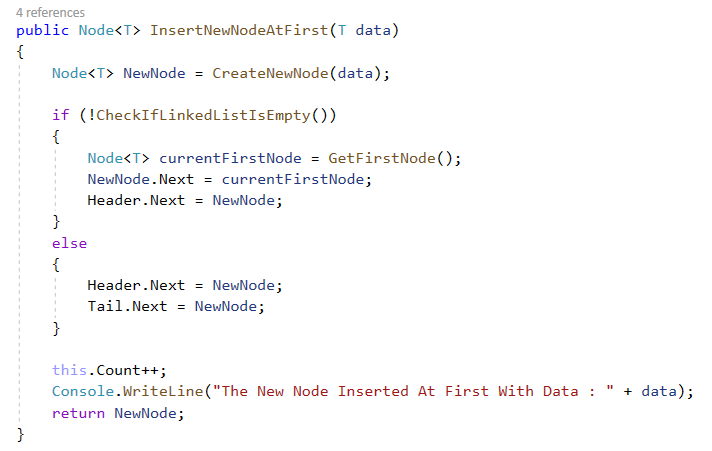


3. Else (When LinkedList Is Empty)

3.1 The Next Of Header Will Point To The New Node

3.2 The Next Of Tail Will Point To The New Node

### - Code :



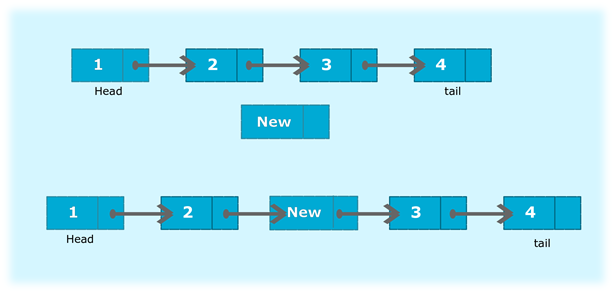
### - Time Complexity :

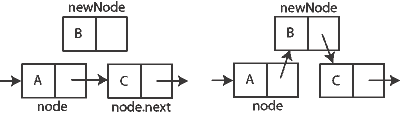
Time Complexity Of Insert At First = O(1)

## 1.2 Insert After Node

The new node will be added after a specific node.

Ex : insert new node after second node :

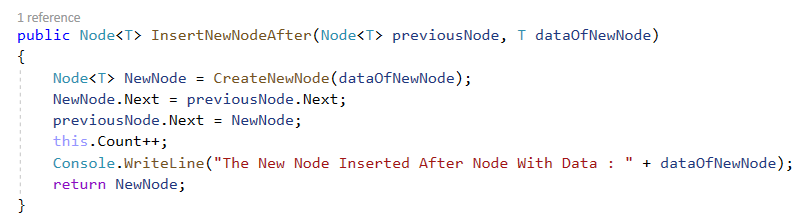




### Steps :

1. Create new node
2. The Next Of New Node Point to The Next Of Previous Node
3. The Next Of Previous Node Point to The New Node

### Code :



### Time Complexity :

Time Complexity Of Insert After Node = O(1)

### 

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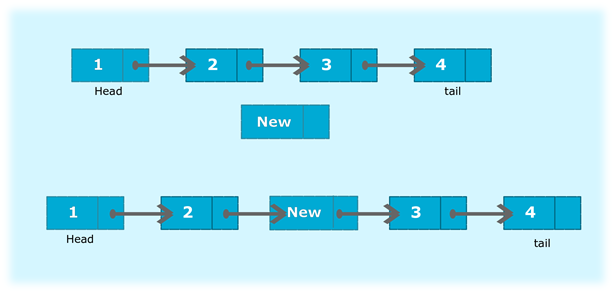
### 

### 

## 1.3 Insert Before Node

The new node will be added before a specific node.

Ex : insert new node before third node :



### Steps :

Inserting a new node before given node operation is not recommended with singly linked list because it has cost of n.

And this Because The Singly Linked List Has Pointer only to the next Node.

We solve this problem By Using Doubly linked List Which Contain Pointers To Next Node And **Also Previous Node**.

We Will Explain The Doubly Linked List In Details Later And We Will See A Lot Of Advantages For Doubly Linked List Over Singly Linked List

### Time Complexity :

Time Complexity Of Insert Before Node = O(n)

## 1.4 Insert At Last

The New Node Will Be The Last Node In The Linked List.

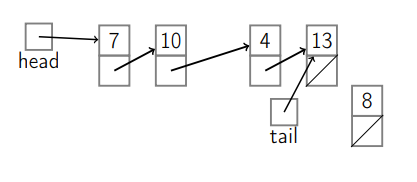
### Steps :

1. If Linked List Is Empty :

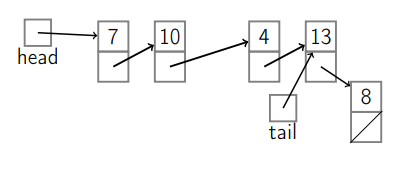
1.1 The New Node Will Insert At First As We See In The Previous Operation

2 Else ( Linked List Is Not Empty )

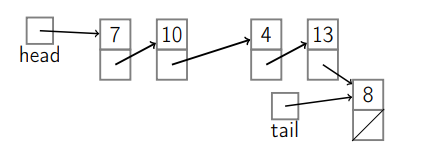
2.1 Create New Node



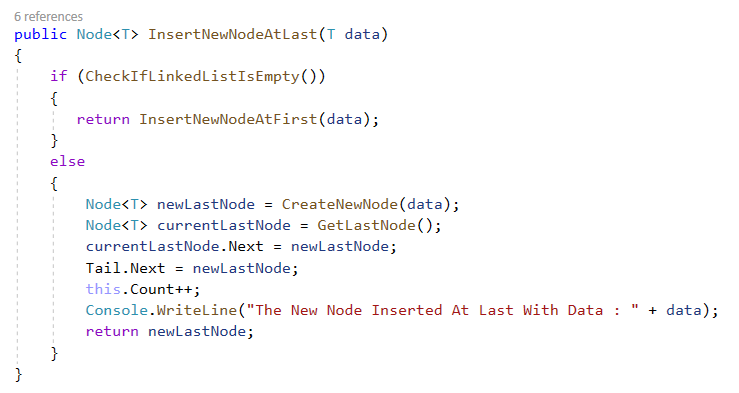
2.2 The Next Of Current Last Node Will Point To The New Node



2.3 The Tail Will Point To The Last Node



### Code :



### Time Complexity

Time Complexity Of Insert At Last = O(1)

Note : We Can Implement Singly LinkedList Without Tail But In This Case The Time Complexity Of Insert At Last = O(n) so we used Tail to prevent this problem

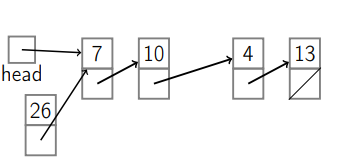
## 2- Delete Operations

## 2.1 Delete The Node At The First

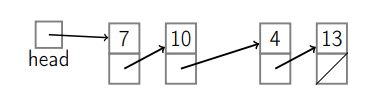
The First Node In The Linked List Will be Deleted And The Current Second Node Will Be The First.

### Steps :

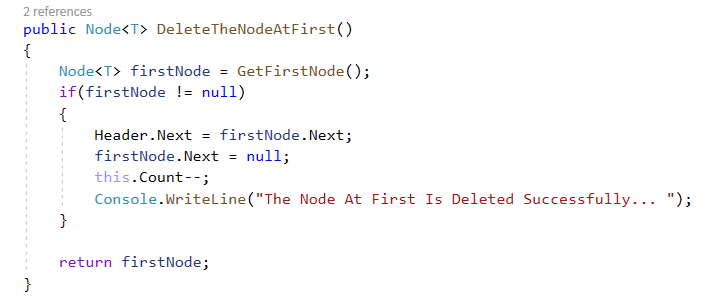
1. The Header Will Point To The Next Of Current First Node.



1. The Next Of Current First Node Will Point To Null



### Code :



### Time Complexity :

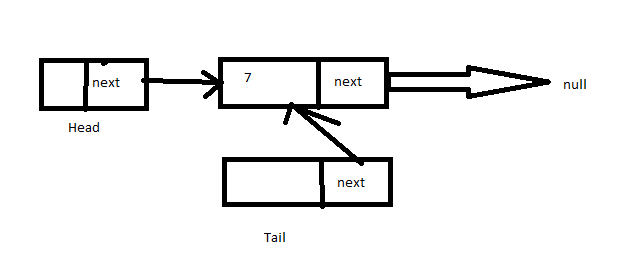
Time Complexity Of Delete Node At First = O(1)

## 2.1 Delete The Node At Last

The Last Node Will Be Deleted And The Node Before The Current Last Node Will Be The New Last Node in The Linked List.

### Steps

1.If Linked List Contain Only One Node

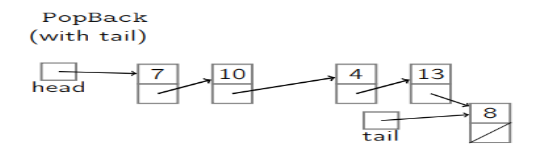


1.1 Delete The Node At First

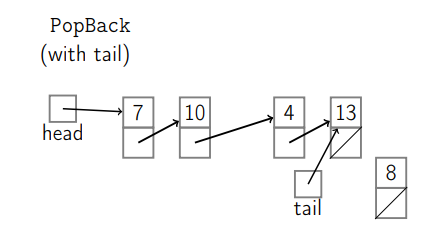
1.2 Header Point To Null

1.3 Tail Point To Null

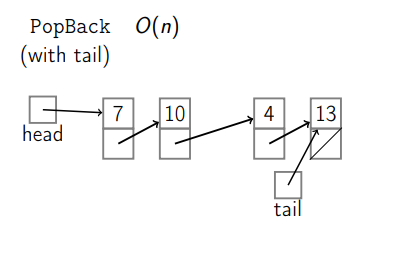
2. Else If Linked List Contain More Than One Node



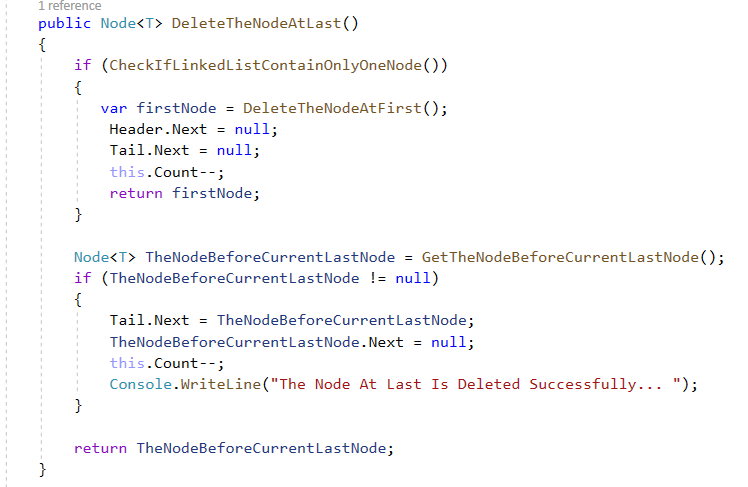
2.1 The Tail Point To The Node That Was Before The Current Last Node.

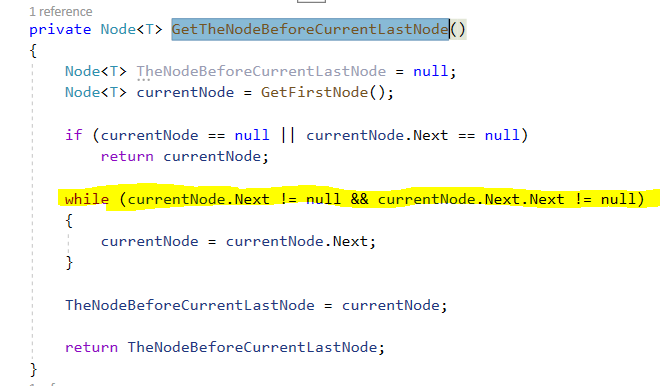


2.2 The Next Of The New Last Node Point To Null



### Code :





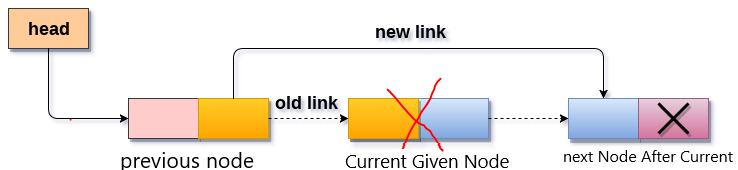
### Time Complexity

Time Complexity Of Delete Node At Last = O(n).

In Singly Linked List , The Node Contain Only Pointer To The Next Node. So We Need To Traverse The Whole Linked List To Get The Node Before Current Last Node.We Solve This Problem By Using Doubly Linked List Which Contain Pointer To The Previous Node, and Get The Node Before The Current Last Node By Just One Step.

## 2.3 Delete Given Node

Delete Current Given Node And Make The Next Of Previous Node Point To The Next Node After The Current Node.



### Steps :

1.If Current Node Is First Node

1.1 Delete Node At First

1.2 RETURN

2.If Current Node Is Last Node

2.1 Delete Node At Last

2.2 RETURN

3. If Current Node != First Node OR Current Node != Last Node

3.1 Get The Previous Node Of The Current Node

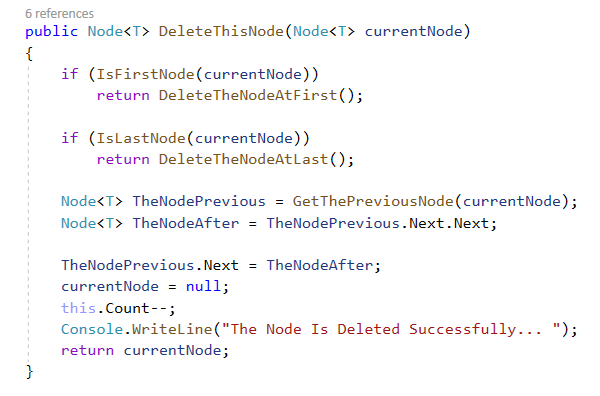
3.2 The Next Of The Previous Node Point To The Node After The Current Node

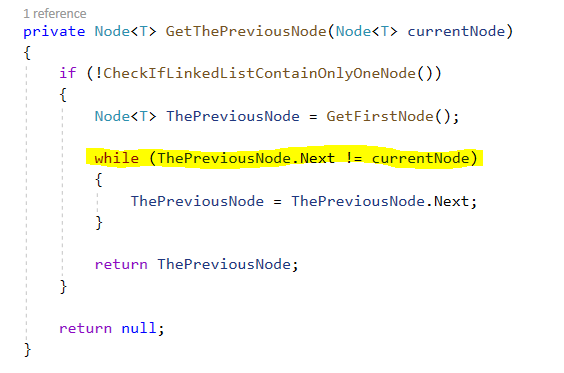
3.3 RETURN

### 

### 

### Code :





### Time Complexity

Time Complexity Of Delete Given Node = O(n).

In Singly Linked List , The Node Contain Only Pointer To The Next Node. So We Need To Traverse The Linked List To Get The Node Before Current Node.We Solve This Problem By Using Doubly Linked List Which Contain Pointer To The Previous Node, and Get The Node Before The Current Node By Just One Step.

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## 2.4 Delete The Node At Specific Position

### Steps :

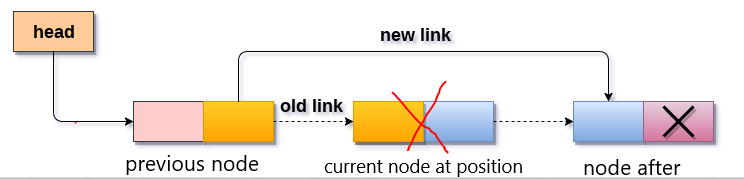
1. If Position == 0

1.1 Delete The Node At First

2. If Position == Linked List Length

2.1 Delete The Node At Last

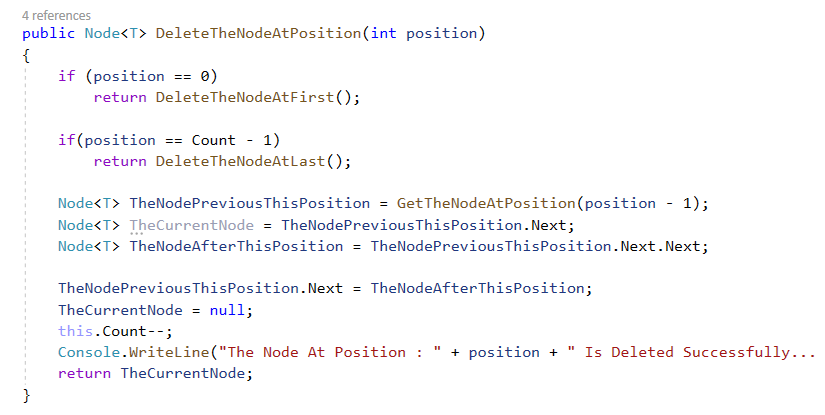
3.If Position != 0 OR Position != Linked List Length

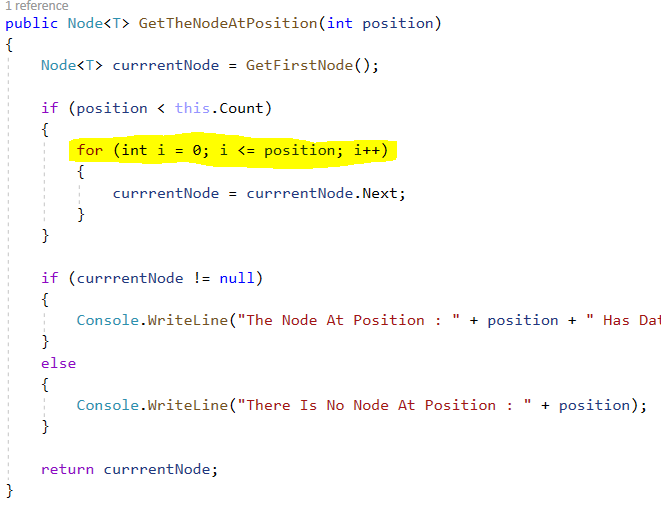


3.1 Get The Previous Node Of This Position

3.2 The Next Of The Previous Node Point To The Node After This Position

### Code :



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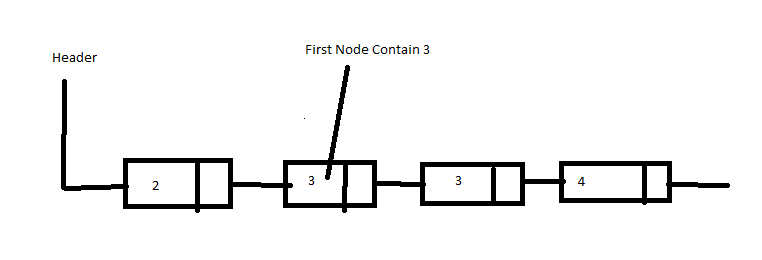
### Time Complexity :

Time Complexity Of Delete Node At Specific Position = O(n).

This Because As We See We Need To Loop The Linked list To Get The Previous Node.

## 3- Search Operation ( Find Node )

Find First Node Contain given Value



### Steps :

1- Define Variable to Store The Current Node

2- Get First Node In The Linked List

3- Current Node = First Node

3- While ( Current Node != null )

3.1 If Data Of Current Node == Value

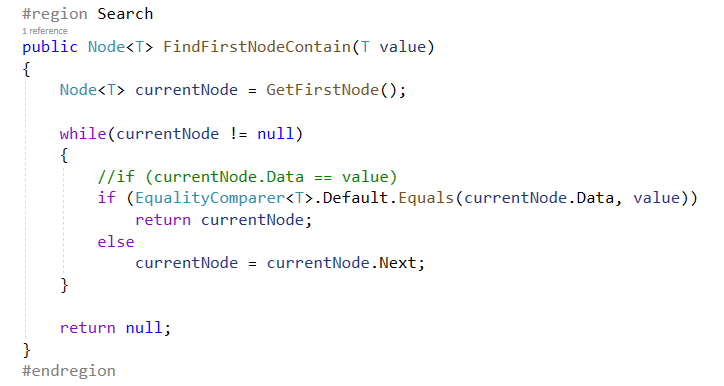
3.2 Return Current Node

3.3 Else ( Data Of Current Node != Value )

3.4 Current Node = Next Node

4- Return Null

### Code :



### Time Complexity :

Time Complexity Of Find First Node Contain Given Value = O(n).

## 4- Read Operations

## 4.1 Get Node At Position

### Steps :

1- Define Variable to Store The Current Node

2- Get First Node In The Linked List

3- If Position < Linked List Length

3.1 Loop (int i = 0; i <= position; i++)

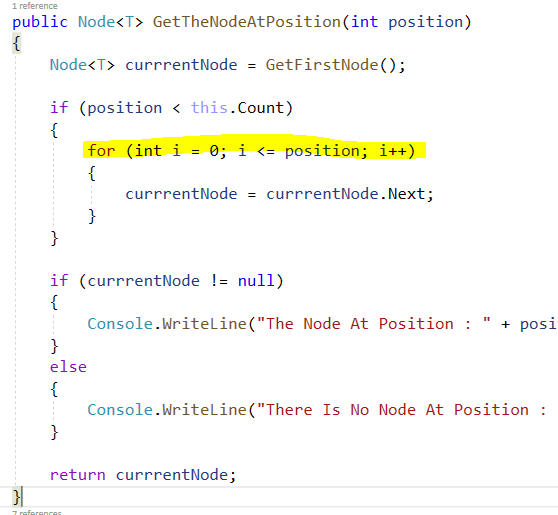
3.2 Current Node = Next Node

3.3 Return Current Node

4. Else Position >= Linked List Length

4.1 Return Null

### Code :



### Time Complexity :

Time Complexity Of Get Node At Position = O(n).

## 4.2 Traverse Linked List

Traverse Linked List To Get Or Print Data Of All Nodes.

### Steps :

1- Define Variable to Store The Current Node

2- Get First Node In The Linked List

3- If Linked List Not Empty

3.1 Print Current Node Data

3.2 Current Node = Next Node

### Code :

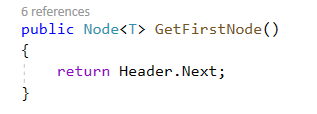


### Time Complexity :

Time Complexity Of Traverse Operation = O(n).

## 4.3 Get Node At First

### Code :

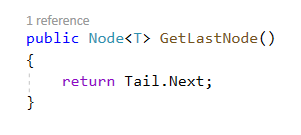


### Time Complexity :

Time Complexity Of Get Node At First = O(1).

## 4.4 Get Node At Last

### Code :



### Time Complexity :

Time Complexity Of Get Node At Last = O(1).